

Student Name:

Student id:

Sect: #:

QUESTION #	1	2	3	4	5	6	TOTAL
MAX POINTS	8	12	10	10	11	12	63
POINTS EARNED							

University of Bahrain

College of Information Technology

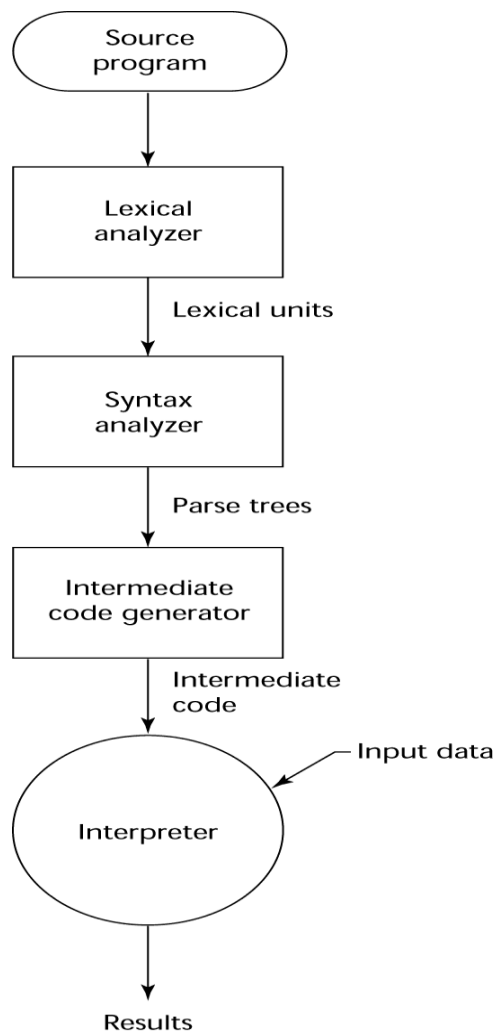
Department of Computer Science

ITCS 332: Organization of Programming Languages FIRST TEST Date: APR 06, 2016

QUESTION ONE:

[8 pts]

Draw the flowchart that illustrates the steps of a hybrid implementation system of any HLL .



QUESTION TWO: Fill in blanks

[12 pts]

- 1) Give two features of imperative languages.
 - Imperative languages are based on assignment and variables.
 - Algorithms in imperative languages consist of detailed ordered steps.
- 2) Give two suggestions to the programmer to reduce the compilation times in C++ language.
 - Use operands of the same type in expressions, if possible.
 - Use explicit casting for types.
 - Use subprograms instead of macros
- 3) Give two suggestions to the programmer to reduce the execution times in C++ language.
 - Use operands of the same type in expressions, if possible.
 - Use explicit casting for types.
 - Use macros instead of subprograms.
- 4) Give two suggestions to the programmer to improve the readability of a language.
 - Use meaningful names.
 - Use familiar language constructs or avoid using unfamiliar ones.
 - Reduce the use of overload operators.
 - Reduce the use of aliases.
- 5) Give two suggestions to the language designer to improve the reliability of a language .
 - Support "natural" ways of expressing algorithms.
 - Exclude aliasing from the language.
 - Include extensive capabilities for exception handling.
 - Remove pointers from the language.
- 6) Give two advantages of using preprocessor macros:
 - Programs with macros are generic and flexible (Macros use Typeless parameters) .
 - Programs with macros are fast in execution.

QUESTION THREE:

[10 pts]

- Convert each of the following BNF rules into ONE equivalent EBNF rule

1) $\langle \text{run} \rangle \rightarrow \langle \text{leg} \rangle \mid \langle \text{run} \rangle \text{"\#" } \langle \text{leg} \rangle \mid \langle \text{run} \rangle \text{"!" } \langle \text{leg} \rangle \mid \langle \text{run} \rangle \text{"\$"} \langle \text{leg} \rangle$

$\langle \text{run} \rangle \rightarrow \langle \text{leg} \rangle \{ (\# \mid ! \mid \$) \langle \text{leg} \rangle \}$

2) $\langle \text{int} \rangle \rightarrow \langle \text{decs} \rangle \# \mid \langle \text{sign} \rangle \langle \text{decs} \rangle \#$
 $\langle \text{decs} \rangle \rightarrow \langle \text{digit} \rangle \mid \langle \text{decs} \rangle \langle \text{digit} \rangle$
 $\langle \text{sign} \rangle \rightarrow + \mid -$

$\langle \text{int} \rangle \rightarrow [(+ \mid -)] \langle \text{digit} \rangle \{ \langle \text{digit} \rangle \} \#$

3) $\langle \text{real} \rangle \rightarrow \langle \text{sign} \rangle \langle \text{decs} \rangle \text{"." } \langle \text{decs} \rangle \text{"E"} \langle \text{decs} \rangle \mid \langle \text{sign} \rangle \langle \text{decs} \rangle \text{"." } \langle \text{decs} \rangle$
 $\mid \langle \text{sign} \rangle \langle \text{decs} \rangle \text{"." } \langle \text{decs} \rangle \text{"E"} \langle \text{sign} \rangle \langle \text{decs} \rangle$
 $\langle \text{decs} \rangle \rightarrow \langle \text{dig} \rangle \mid \langle \text{decs} \rangle \langle \text{dig} \rangle$
 $\langle \text{sign} \rangle \rightarrow \text{"+" } \mid \text{"-" }$

$\langle \text{real} \rangle \rightarrow (+ \mid -) \langle \text{dig} \rangle \{ \langle \text{dig} \rangle \} \text{"." } \langle \text{dig} \rangle \{ \langle \text{dig} \rangle \} [\text{"E"} [(+ \mid -)] \langle \text{dig} \rangle \{ \langle \text{dig} \rangle \}]$

- Fill in blanks**

1) Rewrite the rule: $\langle \text{run} \rangle \rightarrow \langle \text{leg} \rangle \dots$ given above, so that operators # ! \$ associate right to left.

$\langle \text{run} \rangle \rightarrow \langle \text{leg} \rangle \mid \langle \text{leg} \rangle \text{"\#" } \langle \text{run} \rangle \mid \langle \text{leg} \rangle \text{"!" } \langle \text{run} \rangle \mid \langle \text{leg} \rangle \text{"\$"} \langle \text{run} \rangle$

2) In axiomatic semantics, each statement is preceded by a **precondition** and followed by a **postcondition**.

3) Name two issues (examples) in programming languages that cannot be described by BNF rules:

- Type compatibility rules.**
- Identifiers must be declared before referenced.**

QUESTION FOUR:**[10 pts]**

- **Fill in blanks**

- 1) The weakest precondition for each kind of statements can be computed by **an axiom** or by **an inference rule**.
- 2) In denotational semantics, the state changes are defined by **MATHEMATICAL FUNCTIONS**.
In operational semantics, the state changes are defined by **CODED ALGORITHMS**.
- 3) The number of lexemes in a C++ statement: `if(!(xt[j]-9)) xy=ct;` is **16** and the number tokens of is **11**.
- 4) Ambiguities in BNF grammars of arithmetic expression can be solved by introducing **precedence rules** and **assosiativity rule** into its grammar.

- Convert each of the following EBNF rule into equivalent BNF rule(s).

- 1) `<par> → <sent> { ; <sent> }`

```
<par>    → <sent>
          | <sent> ; < par >
```

- 2) `<call> → <id> "(" <parm> { "#" <parm> } ")"`

```
<call>    → <id> "(" <parms> ")"
<parms>   → <parm>
          | <parm> "#" <parms>
```

- 3) `<isa> → <int> { (! | :) <int> }`

```
<isa>     → <int>
          | <int> ! <isa>
          | <int> : <isa>
```

QUESTION FIVE:**[11 pts]**

Carefully study the following grammar and answer the next two questions.

$$\langle E \rangle \rightarrow \langle E \rangle + \langle T \rangle \mid \langle E \rangle - \langle T \rangle \mid \langle T \rangle$$

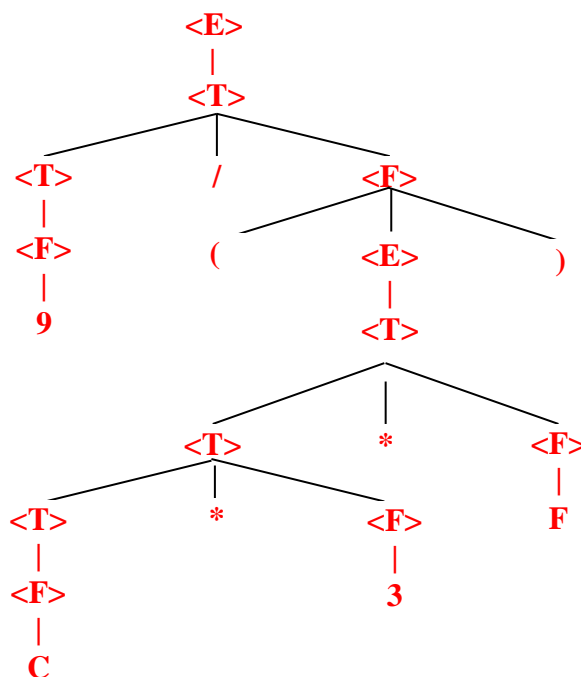
$$\langle T \rangle \rightarrow \langle T \rangle * \langle F \rangle \mid \langle T \rangle / \langle F \rangle \mid \langle F \rangle$$

$$\langle F \rangle \rightarrow (\langle E \rangle) \mid 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9 \mid A \mid B \mid C \mid D \mid E \mid F$$

- Construct the rightmost derivation of the sentence : $5 + (D * A + 9)$

$$\begin{aligned} \langle E \rangle &\Rightarrow \langle E \rangle + \langle T \rangle \\ &\Rightarrow \langle E \rangle + \langle F \rangle \\ &\Rightarrow \langle E \rangle + (\langle E \rangle) \\ &\Rightarrow \langle E \rangle + (\langle E \rangle + \langle T \rangle) \\ &\Rightarrow \langle E \rangle + (\langle E \rangle + \langle F \rangle) \\ &\Rightarrow \langle E \rangle + (\langle E \rangle + 9) \\ &\Rightarrow \langle E \rangle + (\langle T \rangle + 9) \\ &\Rightarrow \langle E \rangle + (\langle T \rangle * \langle F \rangle + 9) \\ &\Rightarrow \langle E \rangle + (\langle T \rangle * A + 9) \\ &\Rightarrow \langle E \rangle + (\langle F \rangle * A + 9) \\ &\Rightarrow \langle E \rangle + (D * A + 9) \\ &\Rightarrow \langle T \rangle + (D * A + 9) \\ &\Rightarrow \langle F \rangle + (D * A + 9) \\ &\Rightarrow 5 + (D * A + 9) \end{aligned}$$

- Construct the parse tree of the sentence : $9 / (C * 3 * F)$



QUESTION SIX:**[12 pts]**

- Construct the BNF rules to define the **cout** statement defined as follows. A **cout** statement is a keyword **cout** or **print** followed by an insertion operator **<<** followed by one or more expressions separated by colons **:** and terminated by **\$**. An expression may be one or more variables or constants. A constant is one or more digits. A digit is **<dig>** $\rightarrow 0|1|\dots|8|9$. Examples of accepted **cout** statements:

```
cout << <var>:345:<var> $
cout << 67 : <var> $
cout << <var> $
```

<cout>	\rightarrow cout << <exprs> \$ print << <exprs> \$
<exprs>	\rightarrow <expr> <expr> : <exprs>
<expr>	\rightarrow <const> <var>
<const>	\rightarrow <dig> <const> <dig>

- Construct the BNF rules to define the **cin** statement defined as follows. A **cin** statement is a keyword **cin** followed by the extraction operator **>>** followed by one or more identifiers separated by colons **:** and terminated by **\$**. An identifier is defined as a letter followed by zero or more digits or letters. Assume letters and digits are already predefined. Examples of accepted **cin** statements:

```
cin >> quiz $
cin >> t : row4 $
cin >> d8 : h : b49d $
```

<cin>	\rightarrow cin >> <ids> \$
<ids>	\rightarrow <id> <id> : <ids>
<id>	\rightarrow <let> <id> <let> <id> <dig>